## I-small submodels of countable models of arithmetic

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## 1970s

Set Theory vs Arithmetic

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- Independence results


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- First appearance: Lascar 1994, Small index property.


## Outline of the talk

- Properties of I-small subsets of $\mathcal{M}$.
- Automorphism group of a countable recursively saturated model of PA and $I$-small submodels.
- Initial self-embeddings of countable models of $I \Sigma_{1}$ and I-small submodels.

Properties of $I$-small subsets of M

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- (Ackermann's membership relation). There exists a $\Delta_{0}$-formula $x E y$ asserting that "the $x$-th bit of the binary expansion of $y$ is 1 ". $a_{\mathrm{E}}$ denotes the set of E -members of $a$ in $\mathcal{M}$.
- $\operatorname{SSy}_{\boldsymbol{I}}(\mathcal{M}):=\left\{X \cap I: X\right.$ is $\Sigma_{1}$-definable in $\left.\mathcal{M}\right\}=\left\{a_{E} \cap I: a \in M\right\}$.
- If $I \subset_{e} \mathcal{M} \subseteq_{e} \mathcal{N}$, then $\operatorname{SSy}_{\boldsymbol{l}}(\mathcal{M})=\operatorname{SSy}_{\boldsymbol{l}}(\mathcal{N})$.


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- $\operatorname{SSy}_{\boldsymbol{\prime}}(\mathcal{M}):=\left\{X \cap I: X\right.$ is $\Sigma_{1}$-definable in $\left.\mathcal{M}\right\}=\left\{a_{E} \cap I: a \in M\right\}$.
- If $I \subset_{e} \mathcal{M} \subseteq_{e} \mathcal{N}$, then $\operatorname{SSy}_{\boldsymbol{l}}(\mathcal{M})=\operatorname{SSy}_{\boldsymbol{l}}(\mathcal{N})$.
$A:=\left\{i \in I: \mathcal{M} \models \neg \operatorname{E}(a)_{i}\right\} \neq \emptyset$ is inside $\operatorname{SSy}_{\boldsymbol{\prime}}(\mathcal{M})$ but not in $\mathrm{SSy}_{\mathrm{I}}\left(\mathcal{M}_{0}\right)$.


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- By Compactness Theorem, there exists some elementary extension $\mathcal{N}$ of $\mathcal{M}$ in which $\mathcal{M}$ is small.


## Which subsets of $M$ are I-small?

## Notation:

- Let $\left\langle\delta_{r}: r \in M\right\rangle$ be a canonical enumeration of all $\Delta_{0}$-formulas within $\mathcal{M}$.
- The predicate $\operatorname{Sat}_{\Delta_{0}}(x)$ is the truth predicate for $\Delta_{0}$-formulas in $\mathcal{M}$, which is $\Delta_{1}$-definable in $\mathcal{M}$.
- For every $r \in M, f_{r}(\bar{x})=y$ denotes the following partial $\Sigma_{1}$-function in $\mathcal{M}$ :

$$
y:=\text { the least element such that } \exists z \operatorname{Sat}_{\Delta_{0}}\left(\delta_{r}(\bar{x}, y, z)\right)
$$

- The notation $\left[f_{r}(\bar{x}) \downarrow\right]$ denotes the $\Sigma_{1}$-formula $\exists z, y \operatorname{Sat}_{\Delta_{0}}\left(\delta_{r}(\bar{x}, y, z)\right.$ ), and $\left[f_{r}(\bar{x}) \downarrow\right]^{<w}$ stands for the formula $\exists z, y<w \operatorname{Sat}_{\Delta_{0}}\left(\delta_{r}(\bar{x}, y, z)\right)$.
- Let $\mathcal{F}$ be the collection of all $\emptyset$-definable partial $\Sigma_{1}$-functions in $\mathcal{M}$.


## Which subsets of $M$ are I-small?

- For every $c \in M$ the subset of $\Sigma_{1}$-definable elements of $\mathcal{M}$ with $c$ as parameter, denoted by $\mathrm{K}^{1}(\mathcal{M} ; c)$ is small in $\mathcal{M}$ :


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It is easy to see that: $\mathrm{K}^{1}(\mathcal{M} ; c)=\left\{f_{n}(c): n \in \mathbb{N}\right.$ and $\left.\mathcal{M} \models\left[f_{n}(c) \downarrow\right]\right\}$.

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\mathrm{K}^{1}(\mathcal{M} ; c)=\left\{f_{n}(c): n \in \mathbb{N} \text { and } \mathcal{M} \models\left[f_{n}(c) \downarrow\right]\right\} \text {. Fix some }
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nonstandard $s \in M$, and let $a \in M$ such that:

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\mathcal{M} \models \forall r<s\binom{\left(\left[f_{r}(c) \downarrow\right] \rightarrow(a)_{r}=f_{r}(c)\right) \wedge}{\left(\neg\left[f_{r}(c) \downarrow\right] \rightarrow(a)_{r}=0\right)} .
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- $\mathrm{K}^{1}(\mathcal{M} ; I)$ is the subset of $\Sigma_{1}$-definable elements of $\mathcal{M}$ with elements of I as parameter.


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Similar to the previous proof, we find some $a \in M$ such that $\mathrm{K}^{1}(\mathcal{M} ; I)=\left\{(a)_{i}: \quad i \in I\right\}$. In order to make the function $(a)_{i}$ an injection, we inductively define the $\Delta_{0}$-function $g$ in $\mathcal{M}$ such that:

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g(0):=(a)_{0}, \text { and }
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$g(x+1):=(a)_{r}$ s.t. $r$ is the least element for which $(a)_{r}$ is not between elements of $\{g(z): z \leq x\}$.

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$g(x+1):=(a)_{r}$ s.t. $r$ is the least element for which $(a)_{r}$ is not between elements of $\{g(z): z \leq x\}$. Then let $h(x):=\mu_{r}\left((a)_{r}=g(x)\right)$. Since $l$ is strong there exists some $e \in M$ s.t. $h(i) \in I$ iff $h(i)<e$ for all $i \in I$.

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So $M_{0}=\left\{g(i): i<i_{0}\right\}$, which is a contradiction. As a result, $g$ li, is a bijection from I onto $\mathrm{K}^{1}(\mathcal{M} ; /)$.

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(2) If $\mathcal{M}$ is a countable recursively saturated model of PA, then: 2-1) (Kossak-Schmerl (1995)). There exists some small recursively saturated elementary submodel $\mathcal{M}_{0}$ of $\mathcal{M}$ which has $2^{N_{0}}$ elementary submodels.

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Let $S$ be a satisfaction class for $\mathcal{M}$ such that $\mathcal{M}^{*}:=(\mathcal{M} ; S)$ is also recursively saturate. Then put $\mathcal{M}_{0}:=\mathrm{K}\left(\mathcal{M}^{*}\right)$

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\text { let } \left.\mathcal{M}_{0}:=\mathrm{K}\left(\mathcal{M}^{*} ; I \cup\{a\}\right) \text { for some } a>1\right) \text {. }
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that $(I ; X) \models \mathrm{PA}^{*}$. So let $(I ; X) \models \mathrm{x}_{0}:=\mu_{x}(\forall y\langle y, x\rangle \notin X)$.
Therefore, $0 \neq x_{0} \notin M_{0}$ but $x_{0}-1 \in M_{0}$.

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## Question.

Is the strongness of I necessary in the previous statements?

I-small submodels and automorphisms of $\mathcal{M}$

## Results about automorphisms of $\mathcal{M}$

Schmerl (in Kaye-Kossak-Kotlarski’s 1991 paper)
Suppose $\mathcal{M}$ is a countable recursively saturated model of $\mathrm{PA}, I$ is a cut of $\mathcal{M}$, and $\mathcal{M}_{0}$ is an I-small elementary submodel of $\mathcal{M}$. Then I is strong in $\mathcal{M}$ iff there exists some automorphism $j$ of $\mathcal{M}$ such that $M_{0}=\operatorname{Fix}(j)$.

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Suppose $\mathcal{M}$ is a countable recursively saturated model of PA. Then:
I) for every small elementary submodel of $\mathcal{M}_{0}$ and every automorphism $j$ of $\mathcal{M}, M_{0} \cap \operatorname{Fix}(j)$ is small in $\mathcal{M}$.

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Schmerl (in Kaye-Kossak-Kotlarski’s 1991 paper)
Suppose $\mathcal{M}$ is a countable recursively saturated model of PA , I is a cut of $\mathcal{M}$, and $\mathcal{M}_{0}$ is an l-small elementary submodel of $\mathcal{M}$. Then I is strong in $\mathcal{M}$ iff there exists some automorphism $j$ of $\mathcal{M}$ such that $M_{0}=\operatorname{Fix}(j)$.

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Suppose $\mathcal{M}$ is a countable recursively saturated model of PA. Then:
I) for every small elementary submodel of $\mathcal{M}_{0}$ and every automorphism $j$ of $\mathcal{M}, M_{0} \cap \operatorname{Fix}(j)$ is small in $\mathcal{M}$.
II) The following are equivalent:

1) $\mathbb{N}$ is strong in $\mathcal{M}$.
2) For every small $\mathcal{M}_{0} \prec \mathcal{M}$ there exists some automorphism j of $\mathcal{M}$ such that $\operatorname{Fix}(j)=M_{0}$.
3) There exists some automorphism $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \subseteq \mathcal{M}(0)$.
4) There exists some automorphism $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \not \equiv \mathcal{M}$.

## Results about automorphisms of $\mathcal{M}$

## Kossak-Kotlarski (1996)

Suppose $\mathcal{M}$ is a countable recursively saturated model of PA, $\mathcal{M}_{0}=$ $\left\{(a)_{n}: n \in \mathbb{N}\right\}$ is a small elementary submodel of $\mathcal{M}$ and $j$ is an automorphism $\mathcal{M}_{0}$. Then there exists an automorphism $\hat{j}$ of $\mathcal{M}$ which extends $j$ iff there exists some $b \in M$ such that $j\left((a)_{n}\right)=(b)_{n}$ for all $n \in \mathbb{N}$, and the same holds for $j^{-1}$.

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## Enayat (2006)

Suppose $\mathcal{M} \models \mathrm{PA}$ is countable, recursively saturated, and $l$ is a strong cut of $\mathcal{M}$. Moreover, let $\mathcal{M}_{0}$ be an l-small elementary submodel of $\mathcal{M}$. Then there exists a group embedding $\Phi$ from $\operatorname{Aut}(\mathbb{Q},<)$ into $\operatorname{Aut}(\mathcal{M})$ such that for every fixed point free automorphism $j$ of $(\mathbb{Q},<)$ it holds that $\operatorname{Fix}(\Phi(j))=M_{0}$.

## I-small submodels and initial self-embeddings of $\mathcal{M}$

## Friedman's Theorem

## Friedman (1973)

Let $\mathcal{M}, \mathcal{N}$ be countable nonstandrd models of PA. The following statements are equivalent:
(1) $\operatorname{SSy}(\mathcal{M})=\operatorname{SSy}(\mathcal{N})$, and $\operatorname{Th}_{\Sigma_{1}}(\mathcal{M}) \subseteq \operatorname{Th}_{\Sigma_{1}}(\mathcal{N})$.
(2) There is an embedding $j: \mathcal{M} \rightarrow \mathcal{N}$ such that $j(\mathrm{M}) \subset_{e} \mathcal{N}$.

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There are many refinements of Friedman's Theorem in the literature. In particular, Ressayre proved a similar result for models of I $\Sigma_{1}$. Moreover, Dimitracopoulos and Paris developed a version of Friedman's Theorem for models of I $\Delta_{0}+$ Exp.

## I-small submodels as fixed point

## B-Enayat (2018)

Suppose $\mathcal{M} \models \mathrm{I} \Sigma_{1}$ is countable and nonstandard and $I$ is a cut of $\mathcal{M}$. Then the following hold:
(1) I is strong in $\mathcal{M}$ and $I \prec_{\Sigma_{1}} \mathcal{M}$, iff there exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j)=1$.
(2) $\mathbb{N}$ is strong in $\mathcal{M}$ iff there exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j)=\mathrm{K}^{1}(\mathcal{M})$.

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B (2022)
Suppose $\mathcal{M} \models \mathrm{I} \Sigma_{1}$ is countable and nonstandard, I is a cut, and $\mathcal{M}_{0}$ is an I-small $\Sigma_{1}$-elementary submodel of $\mathcal{M}$. Then I is strong in $\mathcal{M}$ iff there exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j)=M_{0}$.

## Sketch of proof of left to right:

- $\mathrm{I}^{1}(\mathcal{M} ; \mathrm{X}):=\left\{x: x \leq a\right.$ for some $\left.a \in \mathrm{~K}^{1}(\mathcal{M} ; X)\right\} \prec \Sigma_{0} \mathcal{M} ;$
- $\mathrm{H}^{1}(\mathcal{M} ; X):=\bigcup_{k \in \omega} \mathrm{H}_{k}^{1}(\mathcal{M} ; X)$, where:

$$
\begin{gathered}
\mathrm{H}_{0}^{1}(\mathcal{M} ; X):=\mathrm{I}^{1}(\mathcal{M} ; X) \text {, and } \\
\mathrm{H}_{k+1}^{1}(\mathcal{M} ; X):=\mathrm{I}^{1}\left(\mathcal{M} ; \mathrm{H}_{k}^{1}(\mathcal{M} ; X)\right) .
\end{gathered}
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- $\mathrm{H}^{1}(\mathcal{M} ; X) \prec \Sigma_{1} \mathcal{M}$ and $\mathrm{H}^{1}(\mathcal{M} ; X) \models \mathrm{I} \Sigma_{1}$.


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- $\mathrm{H}^{1}(\mathcal{M} ; X) \prec \Sigma_{1} \mathcal{M}$ and $\mathrm{H}^{1}(\mathcal{M} ; X) \models \mathrm{I} \Sigma_{1}$.
(i) We will construct some proper initial self-embedding $\alpha$ of $\mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$ such that $\operatorname{Fix}(\alpha)=M_{0}$ and $\alpha\left(\mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)\right)<b$ for some $b \in \mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$.


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(ii) By $\mathrm{I} \Sigma_{1}$-version of the Friedman's Theorem, let
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(iii) Finally, put $j:=\beta^{-1} \alpha \beta$.


## Construction of $\alpha$ :

- First by using strong $\Sigma_{1}$-Collection in $\mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$, we will find some $b \in \mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$ such that:

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\mathcal{M} \models\left[f\left((a)_{i}\right) \downarrow\right] \rightarrow\left[f\left((a)_{i}\right) \downarrow\right]^{<b}, \text { for all } f \in \mathcal{F} \text { and all } i \in I .
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- Back and forth: We will build finite functions $\bar{u} \mapsto \bar{v}$ of elements of $\mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$ such that the following properties hold:
- $\mathrm{P}(\bar{u}, \bar{v}, i, f) \equiv\left[f\left(\bar{u},(a)_{i}\right) \downarrow\right] \rightarrow\left[f\left(\bar{v},(a)_{i}\right) \downarrow\right]^{<b}$, for all $f \in \mathcal{F}$ and $i \in I$,


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- $Q(\bar{u}, \bar{v}, i, f) \equiv\left(\begin{array}{c}{\left[f\left(\bar{u},(a)_{i}\right) \downarrow\right] \wedge} \\ {\left[f\left(\bar{v},(a)_{i}\right) \downarrow\right]^{<b} \wedge} \\ f\left(\bar{u},(a)_{i}\right) \notin M_{0}\end{array}\right) \Rightarrow f\left(\bar{u},(a)_{i}\right) \neq f\left(\bar{v},(a)_{i}\right)$, for all
$f \in \mathcal{F}$ and all $i \in I$.


## Construction of $\alpha$ :

Note that $\mathrm{Q}(\bar{u}, \bar{v}, i, f)$ can be written as a $\Pi_{1}$-formula; to be more exact, let:

$$
\begin{aligned}
& R:= \\
& \left\{\langle k, t\rangle \in I: \mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right) \models\binom{\left(\left[f\left(\bar{u},(a)_{i}\right) \downarrow\right] \wedge\left[f_{t}\left(\bar{u},(a)_{k}\right) \downarrow\right]\right) \rightarrow}{f\left(\bar{u},(a)_{i}\right)=f_{t}\left(\bar{u},(a)_{k}\right)}\right\} .
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$R$ is $\Pi_{1}$-definable and so coded in $\mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$.

## Construction of $\alpha$ :

- Forth levels (for making domain of $\alpha$ to be equal to $\mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$ ): Suppose $\bar{u} \mapsto \bar{v}$ is constructed and $m \in \mathrm{H}^{1}\left(\mathcal{M} ; M_{0}\right)$ is arbitrary. So w.l.o.g. we can assume that $m \leq t\left(\bar{u},(a)_{i_{0}}\right)$ for some $t \in \mathcal{F}$ and $i_{0} \in I$.


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For every $s \in H^{1}\left(\mathcal{M} ; M_{0}\right)$ be arbitrary and put:
$p_{s}(y):=$
$\left\{y \leq t\left(\bar{v},(a)_{i_{0}}\right)\right\} \cup\left\{\forall i, i^{\prime}<s\binom{P(\bar{u}, m, \bar{v}, y, i, f) \wedge}{\mathrm{Q}\left(\bar{u}, m, \bar{v}, y, i^{\prime}, f^{\prime}\right)}: f, f^{\prime} \in \mathcal{F}\right\}$.


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Our aim is to find some $s>1$ such that the bounded $\Pi_{1}$-type $p_{s}(y)$ is finitely satisfiable.


## Construction of $\alpha$ :

Define the following function:

$$
\mathrm{G}(x):=\max \{s<b: \mathcal{M} \models \Theta(s, x, \bar{u}, m, \bar{v})\}
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in which $\Theta(s, x, \bar{u}, m, \bar{v})$ is the following $\Delta_{0}$-formula:

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\forall r, r^{\prime}<x \exists y \leq t\left(\bar{v},(a)_{i_{0}}\right) \forall i, i^{\prime}<s\left(\mathrm{P}\left(\bar{u}, m, \bar{v}, y, i, f_{r}\right) \wedge Q\left(\bar{u}, m, \bar{v}, y, i^{\prime}, f_{r^{\prime}}\right)\right)
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Now by strongness of I there exists some $e>1$ such that $\mathrm{G}(i)>1$ iff $\mathrm{G}(i)>e$ for all $i \in I$. We will show that $p_{e}(y)$ is a finitely satisfiable type.

## Construction of $\alpha$ :

## Lemma

For every finite number of elements of $\mathcal{F}$, say $f, f_{n_{1}}, \ldots, f_{n_{k}}$, there exists some s > I such that:

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\mathcal{M} \models \exists y<t\left(\bar{v},(a)_{i_{0}}\right) \forall i, i^{\prime}<s\left(\mathrm{P}(\bar{u}, m, \bar{v}, y, i, f) \wedge \bigwedge_{w=1}^{k} \mathrm{Q}\left(\bar{u}, m, \bar{v}, y, i^{\prime}, f_{n_{w}}\right)\right) .
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First not that by $I \Sigma_{1}$-version of Friedman's Theorem, for all $s>1$ it holds that:
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## The standard cut

## Corollary

Let $\mathcal{M} \models \mathrm{I} \Sigma_{1}$ be countable and nonstandard. T.F.A.E:

1) $\mathbb{N}$ is strong in $\mathcal{M}$.
2) For every small $\mathcal{M}_{0} \prec_{\Sigma_{1}} \mathcal{M}$ there exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j)=M_{0}$.
3) There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \subseteq I^{1}(\mathcal{M})$.

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If $\mathcal{M} \models \mathrm{PA}$ is recursively saturated, then the above statements are equivalent to the following:
4) There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \models B \Sigma_{1}$ and it is isomorphic to no proper initial segments of $\mathcal{M}$.

## The standard cut

## Proof.

$(3) \Rightarrow(1):$ There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \subseteq I^{1}(\mathcal{M})$.

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## The standard cut

## Proof.

(3) $\Rightarrow(1)$ : There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \subseteq I^{1}(\mathcal{M})$. First by strong $\Sigma_{1}$-collection axiom in $\mathcal{M}$, there exists some $b \in M \backslash I^{1}(\mathcal{M})$. Now, suppose $\mathbb{N}$ is not strong in $\mathcal{M}$.

Lemma. Suppose $\mathcal{M} \vDash \mathrm{I} \Sigma_{1}$ in which $\mathbb{N}$ is not a strong cut, and $j$ is a self-embedding of $\mathcal{M}$, then for every element $b \in M$ there exists an element $c \in \operatorname{Fix}(j)$ such that $\operatorname{Th}_{\Sigma_{1}}(\mathcal{M} ; b) \subseteq \operatorname{Th}_{\Sigma_{1}}(\mathcal{M} ; c)$.

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As a result, there exists some $c \in \operatorname{Fix}(j)$ such that $\operatorname{Th}_{\Sigma_{1}}(\mathcal{M} ; b) \subseteq \operatorname{Th}_{\Sigma_{1}}(\mathcal{M} ; c)$. Which is a contradiction.

## The standard cut

## Proof.

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## Proof.

$(2) \Rightarrow(4)$ : Let $\mathcal{M}_{0}$ be the small recursively saturated elementary submodel of $\mathcal{M}$ we talked about in the previous slides. So by (2) there exists some proper initial self-embedding $j$ of $\mathcal{M}$ s.t. $\operatorname{Fix}(j)=M_{0}$. Since $M_{0}$ is small, $\operatorname{SSy}\left(\mathcal{M}_{0}\right) \neq \operatorname{SSy}(\mathcal{M})$, so it cannot be isomorphic to any initial segments of $\mathcal{M}$.
(4) $\Rightarrow$ (1): There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \models B \Sigma_{1}$ and it is isomorphic to no proper initial segments of $\mathcal{M}$.

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$(2) \Rightarrow(4)$ : Let $\mathcal{M}_{0}$ be the small recursively saturated elementary submodel of $\mathcal{M}$ we talked about in the previous slides. So by (2) there exists some proper initial self-embedding $j$ of $\mathcal{M}$ s.t. $\operatorname{Fix}(j)=M_{0}$. Since $M_{0}$ is small, $\operatorname{SSy}\left(\mathcal{M}_{0}\right) \neq \operatorname{SSy}(\mathcal{M})$, so it cannot be isomorphic to any initial segments of $\mathcal{M}$.
(4) $\Rightarrow$ (1): There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \models B \Sigma_{1}$ and it is isomorphic to no proper initial segments of $\mathcal{M}$. If $\mathbb{N}$ is not strong, by the previous Lemma for every $a \in M$, there exists some $b \in \operatorname{Fix}(j)$ such that $\mathbb{N} \cap a_{E}=\mathbb{N} \cap b_{E}$. As a result, $\operatorname{SSy}(\operatorname{Fix}(j))=\operatorname{SSy}(\mathcal{M})$.

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(2) $\Rightarrow$ (4): Let $\mathcal{M}_{0}$ be the small recursively saturated elementary submodel of $\mathcal{M}$ we talked about in the previous slides. So by (2) there exists some proper initial self-embedding $j$ of $\mathcal{M}$ s.t. $\operatorname{Fix}(j)=M_{0}$. Since $M_{0}$ is small, $\operatorname{SSy}\left(\mathcal{M}_{0}\right) \neq \operatorname{SSy}(\mathcal{M})$, so it cannot be isomorphic to any initial segments of $\mathcal{M}$.
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(2) $\Rightarrow$ (4): Let $\mathcal{M}_{0}$ be the small recursively saturated elementary submodel of $\mathcal{M}$ we talked about in the previous slides. So by (2) there exists some proper initial self-embedding $j$ of $\mathcal{M}$ s.t. $\operatorname{Fix}(j)=M_{0}$. Since $M_{0}$ is small, $\operatorname{SSy}\left(\mathcal{M}_{0}\right) \neq \operatorname{SSy}(\mathcal{M})$, so it cannot be isomorphic to any initial segments of $\mathcal{M}$.
(4) $\Rightarrow$ (1): There exists some proper initial self-embedding $j$ of $\mathcal{M}$ such that $\operatorname{Fix}(j) \models B \Sigma_{1}$ and it is isomorphic to no proper initial segments of $\mathcal{M}$. If $\mathbb{N}$ is not strong, by the previous Lemma for every $a \in M$, there exists some $b \in \operatorname{Fix}(j)$ such that $\mathbb{N} \cap a_{E}=\mathbb{N} \cap b_{E}$. As a result, $\operatorname{SSy}(\operatorname{Fix}(j))=\operatorname{SSy}(\mathcal{M})$. Moreover, $\operatorname{Fix}(j) \preceq_{\Sigma_{1}} \mathcal{M}$. So by I $\Delta_{0}+$ Exp-version of the Friedman's Theorem, there exists a proper initial embedding from $\operatorname{Fix}(j)$ into $\mathcal{M}$, which contradicts (4).

## I-small submodels and extendability of initial self-embeddings of $\mathcal{M}$

## Theorem (B 2022)

Suppose $\mathcal{M} \models I \Sigma_{1}$ is countable and nonstandard, I is a strong cut of $\mathcal{M}, \mathcal{M}_{0}$ is an 1 -small $\Sigma_{1}$-elementary submodel of $\mathcal{M}$ such that $M_{0}:=\left\{(a)_{i}: i \in I\right\}$, and $j$ is an initial self-embedding of $\mathcal{M}_{0}$ such that $j(I) \subseteq_{e} \mathcal{M}$. Then the following are equivalent:
(1) $j$ extends to some proper initial self-embedding of $\mathcal{M}$.
(2) $\cdot$ There exists some $b \in M$ such that $\mathcal{M} \models j\left((a)_{i}\right)=(b)_{(i)}$ for all $i \in I$, and

- for every subset $A$ of $M_{0}$ it holds that:

$$
A \in \operatorname{SSy}_{l}(\mathcal{M}) \text { iff } j(A) \in \operatorname{SSy}_{j(l)}(\mathcal{M})
$$

## Thank you!

## Construction of $\alpha$ :

## Lemma

For every finite number of elements of $\mathcal{F}$, say $f, f_{n_{1}}, \ldots, f_{n_{k}}$, there exists some s > I such that:

$$
\mathcal{M} \models \exists y<t\left(\bar{v},(a)_{i_{0}}\right) \forall i<s\left(\mathrm{P}(\bar{u}, m, \bar{v}, y, i, f) \wedge \bigwedge_{w=1}^{k} \mathrm{Q}\left(\bar{u}, m, \bar{v}, y, i, f_{n_{w}}\right)\right) .
$$

## Proof of Lemma:

Suppose not; i.e. there exists the least $k_{0} \in \omega$ for which there exist some $f \in \mathcal{F}$ and $k_{0}$-many elements $f_{n_{1}}, \ldots, f_{n_{k_{0}}}$ of $\mathcal{F}$ such that for all $s>1$ :
(1): $\quad \mathcal{M} \vDash \forall y<t\left(\bar{v},(a)_{i_{0}}\right) \forall i<s \neg\binom{\mathrm{P}(\bar{u}, m, \bar{v}, y, i, f) \wedge}{\bigwedge_{w=1}^{R_{0}} \mathrm{Q}\left(\bar{u}, m, \bar{v}, y, i, f_{n_{w}}\right)}$.

## Construction of $\alpha$ :

To make things a little more clear, by taking a look at $\mathrm{P}(\bar{u}, m, \bar{v}, y, i, f)$ and $\mathrm{Q}\left(\bar{u}, m, \bar{v}, y, i, f_{n_{w}}\right)$, statement (1) states that:
for all $y<t\left(\bar{v},(a)_{i_{0}}\right)$ and all $\epsilon \in M_{0}$ if $\overbrace{[f(\bar{u}, m, \epsilon) \downarrow] \rightarrow[f(\bar{v}, y, \epsilon) \downarrow]}^{\mathrm{P}(\bar{u}, m, \bar{v}, y, \epsilon, f)}$, then there exists some $\xi \in M_{0}$ and some $w=1, \ldots, k_{0}$ s.t.

$$
\neg Q\left(\bar{u}, m, \bar{v}, y, \xi, f_{n}\right)
$$

$$
\overbrace{\left[f_{n_{w}}(\bar{u}, m, \xi) \downarrow\right] \wedge\left[f_{n_{w}}(\bar{v}, y, \xi) \downarrow\right] \wedge f_{n_{w}}(\bar{u}, m, \xi)=f_{n_{w}}(\bar{v}, y, \xi)} .
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$$
\neg Q\left(\bar{u}, m, \overline{\bar{v}}, \boldsymbol{y}, \xi, f_{n_{w}}\right)
$$

$$
\overbrace{\left[n_{n_{w}}(\bar{u}, m, \xi) \downarrow\right] \wedge\left[f_{n_{w}}(\bar{v}, y, \xi) \downarrow\right] \wedge f_{n_{w}}(\bar{u}, m, \xi)=f_{n_{w}}(\bar{v}, y, \xi)} .
$$

By quantifying out $f_{n_{w}}(\bar{u}, m, \xi)$ s from the above statement, it holds that:

There exists some $x$ s.t. for all $y<t\left(\bar{v},(a)_{i_{0}}\right)$ and all $\epsilon \in M_{0}$ if $[f(\bar{u}, m, \epsilon) \downarrow] \rightarrow[f(\bar{v}, y, \epsilon) \downarrow]$, then there exists some $\xi \in M_{0}$ and some $w=1, \ldots, k_{0}$ s.t. $\left[f_{n_{w}}(\bar{u}, m, \xi) \downarrow\right] \wedge\left[f_{n_{w}}(\bar{v}, y, \xi) \downarrow\right] \wedge(x)_{\left\langle n_{w}, \xi\right\rangle}=f_{n_{w}}(\bar{v}, y, \xi)$.

## Construction of $\alpha$ :

Then we separate those subformulas of the above formula which contain parameters $\bar{u}$ and $m$. It turns out that the subsets defined in $\mathrm{H}^{1}(\mathcal{M} ; I)$ with these subformulas can be coded by suitable elements of $M_{0}$. As a result, we will have a $\Sigma_{1}$-formula whose parameters are only $\bar{v}$ and some elements from $M_{0}$, say $(a)_{i_{1}}$ and $(a)_{i_{2}}$, which serve as codes the aforementioned subsets of $\mathrm{H}^{1}(\mathcal{M} ; I)$; i.e. it holds that:

There exists some $x$ s.t. for all $y<t\left(\bar{v},(a)_{i_{0}}\right)$ and all $\epsilon \mathrm{E}(a)_{i_{1}}$ if $[f(\bar{v}, y, \epsilon) \downarrow]$, then there exists some $\xi \mathrm{E}(a)_{i_{2}}$ and some $w=1, \ldots, k_{0}$ s.t.

$$
\left[f_{n_{w}}(\bar{v}, y, \xi) \downarrow\right] \wedge(x)_{<n_{w}, \xi>}=f_{n_{w}}(\bar{v}, y, \xi) .
$$

## Construction of $\alpha$ :

## Let:

$$
g(\bar{v}):=
$$

the smallest $x$ s.t. for all $y<t\left(\bar{v},(a)_{i_{0}}\right)$ and all $\in \mathrm{E}(a)_{i_{1}}$, if $[f(\bar{v}, y, \epsilon) \downarrow]$, then there exists some $\xi \mathrm{E}(a)_{i_{2}}$ and some $w=1, \ldots, k_{0}$ s.t.

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$$

Moreover, we define:

$$
<o(\bar{v}, y), h(\bar{v}, y)>:=
$$

the smallest $<n_{w}, \xi>$ s.t $1 \leq w \leq k_{0}$ and $\xi \mathrm{E}(a)_{i_{2}}$ and $\epsilon \mathrm{E}(a)_{i_{1}}$, if $[f(\bar{v}, y, \epsilon) \downarrow]$, then $\left[f_{n_{w}}(\bar{v}, y, \xi) \downarrow\right] \wedge(g(\bar{v}))_{<n_{w}, \xi>}=f_{n_{w}}(\bar{v}, y, \xi)$.

## Construction of $\alpha$ :

So it holds that:

$$
\mathcal{M} \models \forall y<t\left(\bar{v},(a)_{i_{0}}\right) \forall \epsilon \mathrm{E}(a)_{i_{1}}\binom{[f(\bar{v}, y, \epsilon) \downarrow] \rightarrow}{[<o(\bar{v}, y), h(\bar{v}, y)>\downarrow]} .
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Then by induction hypothesis:

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If $k_{0}>1$ :
$\mathcal{M} \models \forall y<t\left(\bar{v},(a)_{i_{0}}\right) \forall i<s \neg\binom{P\left(\bar{u}, m, \bar{v}, y, i, f^{\prime}\right) \wedge}{\bigwedge_{w=2}^{k_{0}} \mathrm{Q}\left(\bar{u}, m, \bar{v}, y, i, f_{n_{w}}\right)}$; in which $f^{\prime}$
is:
$f^{\prime}(\diamond, y)=\Leftrightarrow f(\diamond, y)=\wedge[<o(\diamond, y), h(\diamond, y)>]$.

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is:
$f^{\prime}(\diamond, y)=\Leftrightarrow f(\diamond, y)=\wedge[<0(\diamond, y), h(\diamond, y)>]$.
But this contradicts the minimality of $k_{0}$.

## Construction of $\alpha$ :

If $k_{0}=1:$

- $\mathcal{M}$ thinks that the cardinality of

$$
A:=\left\{h(\bar{u}, y): \mathcal{M} \models\left(y<t\left(\bar{u},(a)_{i_{0}}\right) \wedge[h(\bar{u}, y) \downarrow]\right)\right\} \text { is at most }
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- By using the previous statements, we can build a coded 1-1 function $F \in M$ whose domain contains $I$ and $F(I) \subset A$.


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- By using the previous statements, we can build a coded 1-1 function $F \in M$ whose domain contains $I$ and $F(I) \subset A$.

So again a contradiction is achieved by $\Sigma_{1}$-Pigeonhole Principle.

